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A New Type of Shifted Method for the Stein-Like Matrix Equation X+M(X) = CShuai Liu1*

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Abstract	Review Article

In this paper, we presented the shifted global full orthogonalization method (SGI-FOM(m)) and shifted global generalized minimum residual method (SGI-GMRES(m)) to solve the Stein-like matrix equation X + M(X) = C with M(X) = AXB

M (X) = $AX^{\cdot}B$, M (X) = $A\overline{X}B$ or M (X) = $AX^{H}B$, respectively, where X is an unknown matrix to be solved. By utilizing the collinearity of the residuals of the seed system M (X) = C and the shifted system X + M(X) = C, the solution of the Stein-like matrix equation can be derived. Some results are also given. Numerical experiments show its efficiency.

Keywords: FOM method; GMRES method; Stein-like matrix equation; collinearity.

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1. INTRODUCTION

In this paper, we consider the following Stein-like matrix equation:

X + M(X) = C(1)

With M (X) = AXB, M (X) = AX'B, M (X) = $A\overline{X}B$ or M (X) = $AX^{H}B$ and A, B, C being known matrices, If M (X) = AXB, Eq. (1) becomes the standard Stein matrix equation which plays an important role in the fields of discrete-time system [1], statistics [2, 3], probability [4] and spectral analysis [5]. For further details on this topic, we refer to see [6] and the references therein. Under the condition that $\lambda_i(A)\lambda_i(B) \neq -1$ for all i, j, the solution X exists and is unique for any C (e.g., [7]). Necessary and sufficient conditions for the existence and uniqueness of the solution X were obtained in [8]. We also refer the reader to [17] for a solution expressed in terms of the companion form of the given matrices, and to [18] for related computational considerations. Recently, the solvability conditions of ⊤-Stein equation, given by $X = AX \cdot B + C$, have been analyzed in [11-13]. In [14, 15], some explicit solutions of the matrix equation $X - A\overline{X}B = C$ have been established. More generally, Zhou *et al.*, [16] discussed the solvability, existence of unique solution, closed-form solution and numerical solution of matrix equation X = Af(X)B + C with $f(X) = X^{\bullet}, f(X) = \overline{X}$ or $f(X) = X^{H}$.

Recently, many iterative methods for solving Eq. (1) have received great attention. Some Arnoldi-type matrix Krylov subspace methods were presented to solve the Stein matrix equation X = AXB + C in [17-23]. By Smith-type iterative methods, Zhou et al., [24, 25] gained the numerical solution under some restrictions on coefficient matrices. While in [26, 27], Wu et al., obtained the iterative solution of Eq. (1) by extending the CGNE method [28] and the gradient-based iterative (GI) method [29]. The CGNE method denotes the conjugate gradient (CG) method applied to the normal equation of the linear system Ax = b. The basic idea of the GI method is to regard the unknown matrix as the parameter of a system

to be identified. However, the extended CGNE method and the extended GI method have shown slow convergence rate in numerical experiments, especially when dealing with the complex matrix equations.

Li and Huang [30] make full use of the shifted structure of X + AXB = C to solve the matrix equations, inspired by this, we consider using a seed system to solve the shifted system. In [31], Frommer and Glassner solved the shifted linear systems by forcing the residual of the shifted system to be collinear to the residual of the seed system. To my knowledge, no one has yet solve the Stein-like matrix equations by forcing the residual of the shifted system to be collinear to the residual of the seed system. In the present work, we are interested in using shifted global full orthogonalization method (SGI-FOM(m)) and shifted global generalized minimum residual method (SGI-GMRES(m)) to solve Eq. (1). To derive the new method, we define the seed system as follows:

 $\mathbf{M}\left(X\right) = C \tag{2}$

With M (X) = AXB, M (X) = AX' B, M (X) = $A\overline{X}B$ or M (X) = $AX^{H}B$. Then, the system X + M (X) = C may be called the shifted system generated by the seed system (2).

We give the following notations to be used later. The notations tr(A) and A^T denote the trace and transpose of a matrix A, respectively. For two matrices A, B $\in \mathbb{R}^{n \times n}$, A \otimes B is their Kronecker product and $\langle A, B \rangle_F = tr(A^TB)$ is the inner product which generates the Frobenius norm as $||A||_F = \sqrt{\langle A, A \rangle_F} = \sqrt{tr(A^A)}$. The matrices A, B are called orthogonal if $\langle A, A \rangle_F = 0$.

The rest of this paper is organized as follows. In Section 2, we review the generalized global arnoldi process. In Section 3, the shifted Gl-FOM(m) method and shifted Gl-GMRES(m) method are given. Some numerical examples are given to illustrate the efficiency of the proposed method in Section 4. Finally, the paper is ended with a brief conclusion in Section 5.

2. Generalized global Arnoldi process

In this Section, we review the generalized global arnoldi process in [30].

Definition 2.1 ([32]) Let $A = [A_1, A_2, ..., A_p]$ and $B = [B_1, B_2, ..., B_l]$ be matrices of dimension $n \times ps$ and $n \times ls$, respectively, where $A_i, B_j \in \mathbb{R}^{n \times s}$ (i = 1, ..., p; j = 1, ..., l). Then the $p \times l$ matrix $A^{\top} \circ_F B$ is defined by

$$A^{\bullet} \diamond_{F} B = \begin{pmatrix} \langle A_{1}, B_{1} \rangle_{F} & \langle A_{1}, B_{2} \rangle_{F} & \cdots & \langle A_{1}, B_{l} \rangle_{F} \\ \langle A_{2}, B_{1} \rangle_{F} & \langle A_{2}, B_{2} \rangle_{F} & \cdots & \langle A_{2}, B_{l} \rangle_{F} \\ \vdots & \vdots & \vdots & \vdots \\ \langle A_{p}, B_{1} \rangle_{F} & \langle A_{p}, B_{2} \rangle_{F} & \cdots & \langle A_{p}, B_{l} \rangle_{F} \end{pmatrix}$$

For a given matrix $V \in \mathbb{R}^{n \times p}$, we define the matrix Krylov subspace as follows:

 $\mathbf{K}_{m}(\mathbf{M}, V) = span\{V, \mathbf{M}(V), ..., \mathbf{M}^{m-1}(V)\}$ (3)

Here $M^{i}(V)$ is defined recursively as $M(M^{i-1}(V))$ and $M^{0}(V) = V$. We describe the generalized global Arnoldi process global Arnoldi process as follows.

Algorithm 2.1 Generalized global Arnoldi process
1. Set $\beta = V F$, $V_1 = V/\beta$.
2. For <i>j</i> = 1, 2,, <i>m</i> , do
3. $W = \mathbf{M}(V_j);$
4. For <i>i</i> = 1, 2,, <i>j</i> , do
5. $h_{ij} = \langle W, V_i \rangle F$;
6. $W = W - h_{ij} V_j$;
7. Enddo;
8. $h_{j+1,j} = W F$, If $h_{j+1,j} = 0$, Stop;
9. $V_{j+1} = W/h_{j+1,j};$
10. Enddo.

Let $V_m = [V_1, V_2, ..., V_m]$ and H_m be the $(m + 1) \times m$ upper Hessenberg matrix whose nonzero entries h_{ij} are

computed by Algorithm 2.1. Suppose that H_m is the m×m matrix obtained from H_m by deleting its last row. It is not difficult to verify that the global Arnoldi process produces an F-orthonormal basis V_1 , V_2 , ..., V_m for the matrix Krylov subspace $K_m(M,V)$, i.e., $V_m \diamond_F V_m = I_m$.

Theorem 2.1 ([30]) Let V_m be defined as above. Then we have

$$\left\| \mathbf{V}_{m} \left(\boldsymbol{\alpha} \otimes \boldsymbol{I}_{p} \right) \right\|_{F} = \left\| \boldsymbol{\alpha} \right\|_{2} \tag{4}$$

Where $\alpha = (\alpha_1, \alpha_2, ..., \alpha_m)$ • is a real vector of \mathbb{R}^m .

Theorem 2.2 ([30]) Let V_m , H_m and H_m , be defined as above. Then the following relations hold

$$\left[\mathbf{M}\left(V_{1}\right),\mathbf{M}\left(V_{2}\right),\ldots,\mathbf{M}\left(V_{m}\right)\right]=V_{m+1}\left(\hat{H}_{m}\otimes I_{p}\right)\ldots\ldots\ldots\ldots\ldots\ldots\ldots\ldots\ldots\ldots(5)$$

Where $e_m^{\bullet} = [0, ..., 0, 1]_{1 \times m}$.

3. Shifted GI-FOM(m) and Shifted GI-GMRES(m) algorithms

In this section, we propose the Shifted Gl-FOM(m) and Shifted Gl-GMRES(m) algorithms for solving Eq. (1).

3.1 Shifted Gl-FOM(m) algorithm

Let $X_0 \in \mathbb{R}^{n \times p}$ be an initial guess and $R_0 = C - M(X_0)$ be the corresponding residual. At the *m*th iterate, the GI-FOM algorithm computes the new approximate solution X_m such that

 $X_m \in X_0 + K(M, R_0), \dots$ (8)

and

$$R_m = C - M(X_m) \perp_F K_m(M, R_0).$$
 (9)

The symbol \perp_F denotes the orthogonality with respect to the inner product $\langle \cdot, \cdot \rangle_F$. Considering the *F*-orthonormal basis $V_m = [V_1, V_2, \dots, V_m]$ of $K_m(M, R_0)$ constructed by the generalized global Arnoldi process, we have:

where the vector $y_m = (y_m^{(1)}, y_m^{(2)}, \dots, y_m^{(m)})^{\bullet} \in \mathbb{R}^m$ is obtained by the orthogonality condition (9). Substituting (10) to (9), by (6), we get

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$$\begin{aligned} R_m &= C - \mathbf{M} \left(X_0 + \mathbf{V}_m (y_m \otimes I_p) \right) \\ &= R_0 - \mathbf{M} \left(\mathbf{V}_m (y_m \otimes I_p) \right) \\ &= \beta V_1 - \left[\mathbf{M} \left(V_1 \right), \mathbf{M} \left(V_2 \right), \dots \mathbf{M} \left(V_m \right) \right] (y_m \otimes I_p) \\ &= \mathbf{V}_m (\beta e_1 \otimes I_p) - \mathbf{V}_m (H_m \otimes I_p) (y_m \otimes I_p) - h_{m+1,m} \mathbf{V}_{m+1} (\mathbf{e}_m^* \otimes I_p) (y_m \otimes I_p) \\ &= \mathbf{V}_m ((\beta e_1 - H_m y_m) \otimes I_p) - h_{m+1,m} \mathbf{V}_{m+1} (\mathbf{e}_m^* y_m \otimes I_p). \end{aligned}$$
(11)

According to the orthogonality condition (9), we have

$$0 = \mathbf{V}_{m}^{\bullet} \diamond_{F} \mathbf{R}_{m}$$

= $\mathbf{V}_{m}^{\bullet} \diamond_{F} \left(\mathbf{V}_{m} ((\beta e_{1} - H_{m} y_{m}) \otimes I_{p}) - h_{m+1,m} \mathbf{V}_{m+1} (e_{m}^{\bullet} y_{m} \otimes I_{p}) \right)$
= $\mathbf{V}_{m}^{\bullet} \diamond_{F} \mathbf{V}_{m} (\beta e_{1} - H_{m} y_{m}) - h_{m+1,m} (\mathbf{V}_{m}^{\bullet} \diamond_{F} \mathbf{V}_{m+1}) (e_{m}^{\bullet} y_{m}).$

Meanwhile, by (5), R_m can also be represented as: $P - C - M (X + V (y \otimes I))$

$$R_{m} = C - M (X_{0} + V_{m}(y_{m} \otimes I_{p}))$$

$$= R_{0} - M (V_{m}(y_{m} \otimes I_{p}))$$

$$= \beta V_{1} - \left[M (V_{1}), M (V_{2}), \dots M (V_{m})\right](y_{m} \otimes I_{p})$$

$$= V_{m+1}(\beta \hat{e}_{1} \otimes I_{p}) - V_{m+1}((\hat{H}_{m}y_{m}) \otimes I_{p})$$

$$= V_{m+1}((\beta \hat{e}_{1} - \hat{H}_{m}y_{m}) \otimes I_{p})$$

$$= V_{m+1}(\omega_{m+1} \otimes I_{p})$$
(13)

Where $\hat{e}_1 = [0, ..., 0, 1]_{1 \times (m+1)}$, and $\omega_{m+1} = \beta \hat{e}_1 - \hat{H}_m y_m$.

For the shifted system (1), we have

$$\overline{X}_{m} = \overline{X}_{0} + \sum_{i=1}^{m} \overline{y}_{m}^{(i)} V_{i} = \overline{X}_{0} + V_{m} (\overline{y}_{m} \otimes I_{p}) \dots (14)$$

$$\overline{R}_{m} = C - \overline{X}_{m} - \mathbf{M} (\overline{X}_{0} + V_{m} (\overline{y}_{m} \otimes I_{p})) \dots (15)$$

Now assume the initial residuals are colinear,

Then we require that

$$\overline{R}_m = \beta_m R_m, \beta_m \in R \quad \dots \quad (17)$$

According to equations (7), (13), (15) and (16), the colinearity condition (17) can be written as:

Multiply equation (18) from left by V_{m+1}^{\bullet} , we have

$$\beta_0 \beta \hat{e}_1 - H_m \bar{y}_m = \beta_m \omega_{m+1}$$

Which leads to a linear system for the unknown \overline{y}_m and β_m ,

$$\left(\bar{H}_{m} \quad \omega_{m+1}\right) \begin{pmatrix} \bar{y}_{m} \\ \beta_{m} \end{pmatrix} = \beta_{0} \beta \hat{e}_{1}.$$
 (19)

The shifted GI-FOM algorithm requires the storage of V_m . As m increases it becomes expensive. To overcome the storage problem of the shifted GI-FOM algorithm, one can restart the algorithm every *m* iterations. Then the restarted and shifted GI-FOM algorithm denoted by SGI-FOM(m) is given as follows.

Algorithm 3.1 SGI-FOM(m)	
1. Choosing X_0 , and a tolerance ε , set $\overline{X}_0 = X_0$, set $\beta_0 = 1$;	
2. Compute $R_0 = C - M(X_0)$ and $\overline{R}_0 = C - \overline{X}_0 - M(\overline{X}_0)$, set V=R ₀ ;	
3. Construct the <i>F</i> -orthonormal basis V_1 , V_2 ,, V_m by Algorithm 2.1.	
4. Determine y_m as the solution of $H_m y_m = \beta e_1$;	
5. Determine \overline{y}_m and β_m as the solution of $(\overline{H}_m \omega_{m+1}) \begin{pmatrix} \overline{y}_m \\ \beta_m \end{pmatrix} = \beta_0 \beta \hat{e}_1;$	
6. Compute $X_m = X_0 + V_m(y_m \otimes I_p)$ and $R_m = C - M(X_m)$;	
7. Compute $\overline{X}_m = \overline{X}_0 + V_m(\overline{y}_m \otimes I_p)$ and $\overline{R}_m = C - \overline{X}_m - M(\overline{X}_m)$;	
8. If $\ \overline{R}_m\ _F < \varepsilon$, then stop; else, set $X_0 = X_m, R_0 = R_m, V = R_0, \overline{X}_0 = \overline{X}_m$	
$\overline{R}_{_{0}}=\overline{R}_{_{m}},eta_{_{0}}=eta_{_{m}}$, go to 3.	

3.2 Shifted Gl-GMRES(m) algorithm

Like the SGI-FOM(m) algorithm, the *m*th iterate X_m of the GI-GMRES algorithm belongs to the affine matrix Krylov subspace $X_0 + K(M, R_0)$. On the other hand, in the GL-GMRES algorithm, the vec y_m in (10) is obtained by imposing the following orthogonality condition:

$$R_{m} = C - M(X_{m}) \perp_{F} K_{m}(M, M(R_{0})).$$
(20)

From (13), we have

As $V_{m+1}^H \diamond_F V_{m+1} = I_{m+1}$, we have $\|R_m\|_F = \|\beta \hat{e}_1 - \hat{H}_m y_m\|_2$.

The orthogonality condition (20) shows that y_m is the solution of the following minimization problem

$$\min_{\mathbf{y}\in R^m}\left\|\boldsymbol{\beta}\hat{\boldsymbol{e}}_1-\hat{\boldsymbol{H}}_m\boldsymbol{y}_m\right\|_2.$$

Consider the QR factorization of the (m+1) × m matrix \hat{H}_m , i.e., $\hat{Q}_m \hat{H}_m = \hat{R}_m$, where \hat{Q}_m is a product of Givens rotations and \hat{R}_m is a upper triangular. Assume that

$$\hat{g}_m = \beta \hat{Q}_m \hat{e}_1,$$

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And R_m denotes the m × m matrix obtained from \hat{R}_m by deleting its last row and g_m is the m-dimensional vector obtained from \hat{g}_m by deleting its last component. Straightforward computations show that $y_m = R_m^{-1} g_m$.

For the shifted system (1), equations (15) to (18) hold for Gl-GMRES algorithm, so we can obtain \overline{y}_m and β_m by solving the linear system:

$$\begin{pmatrix} \overline{H}_m & \omega_{m+1} \end{pmatrix} \begin{pmatrix} \overline{y}_m \\ \beta_m \end{pmatrix} = \beta_0 \beta \hat{e}_1.$$

Like the SGI-FOM algorithm, to remedy the storage problem, the restarted and shifted GI-GMRES algorithm denoted by SGI-GMRES(m) is given as follows.

Algorithm 3.2 SGI-GMRES(m)	
1. Choosing X_0 , and a tolerance ε , set $\overline{X}_0 = X_0$, set $\beta_0 = 1$;	
2. Compute $R_0 = C - M(X_0)$ and $\overline{R}_0 = C - \overline{X}_0 - M(\overline{X}_0)$, set V=R ₀ ;	
3. Construct the <i>F</i> -orthonormal basis V_1 , V_2 ,, V_m by Algorithm 2.1.	
4. Determine y_m as the solution of following least square problem	
$\min_{\boldsymbol{y}\in \boldsymbol{R}^{m}}\left\ \boldsymbol{\beta}\hat{\boldsymbol{e}}_{1}-\hat{\boldsymbol{H}}_{m}\boldsymbol{y}_{m}\right\ _{2};$	
5. Determine \overline{y}_m and β_m as the solution of $(\overline{H}_m \omega_{m+1}) \begin{pmatrix} \overline{y}_m \\ \beta_m \end{pmatrix} = \beta_0 \beta \hat{e}_1;$	
6. Compute $X_m = X_0 + V_m(y_m \otimes I_p)$ and $R_m = C - M(X_m)$;	
7. Compute $\overline{X}_m = \overline{X}_0 + V_m(\overline{y}_m \otimes I_p)$ and $\overline{R}_m = C - \overline{X}_m - \mathbf{M}(\overline{X}_m)$;	
8. If $\ \overline{R}_m\ _F < \varepsilon$, then stop; else, set $X_0 = X_m, R_0 = R_m, V = R_0, \overline{X}_0 = \overline{X}_m$	
$\overline{R}_0 = \overline{R}_m, eta_0 = eta_m$, go to 3.	

4. Numerical examples

In this section, we give some numerical examples to illustrate the effectiveness of the SGI-FOM(m) algorithm and SGI-GMRES(m) algorithm for solving Eq. (1). Also, we compare our new methods with the extended CGNE and GI methods from [25, 26]. The experiments are performed in MATLAB. All tests start from the zero matrix with suitable size and stop as soon as $|| R_m ||_F \le 10^{-9}$.

Example 4.1 We consider the Stein matrix equation X + AXB = C

Where A = [diag([1:100]) + diag(1*ones(299,1),1)] and B = [diag([1:100]) + diag(-1*ones(299,1),-1)].

Here C is chosen such that X = tridiag(-1, 0, 1) is the exact solution of the system.

The numerical comparison results are given in Fig. 1. As it is seen, the SGI-FOM (10) and SGI-GMRES (10) algorithms have almost the same convergence behavior and require fewer iteration steps than the the extended CGNE and GI methods for achieving the prescribed stopping criterion.



Figure 1: Comparison results of Example 4.1

Example 4.2 We consider the • -Stein matrix equation $X + AX \cdot B = C$

where $A = B = tridiag(-1, 4, -1) \in \mathbb{R}^{n \times n}$. Here C is chosen such that the exact solution of the system is X = tridiag(-1, 0, 1).

When n = 200, the numerical comparison results are depicted in Fig. 2. We see from Table 2 that the SGI-FOM(m) and SGI-GMRES(10) algorithms have significant advantages in terms of convergence speed compared to extended CGNE and GI methods.



Figure 2: Comparison results of Example 4.2

Example 4.3 We consider the Stein-like matrix equation

 $X + A\overline{X}B = C$ Where $A = tridiag(-i, 3, i) \in C^{n \times n}$ and $B = tridiag(-i, 0, i) \in C^{s \times s}$. Here C is chosen such that the exact solution of the system is $X = (x_{ii})_{n \times s}$ with $x_{ii} = 1 + i$.

When n = 200 and s = 100, the numerical comparison results are given in Fig.3. As it is seen, the SGI-FOM(10) and SGI-GMRES(10) algorithms require fewer iteration steps than the the extended CGNE and GI methods for achieving the prescribed stopping criterion.



Figure 3: Comparison results of Example 4.3

Example 4.4 We consider the Stein-like matrix equation

 $X + AX^{H}B = C$

where the matrices A, B, and C are generated by the following MATLAB codes.

rand('state',0)

 $\begin{array}{l} A=diag(10+diag(rand(n)))+triu(rand(n,n),1)*i\\ B=diag(10+diag(rand(n)))+tril(rand(n,n),1)*i\\ C=rand(n,n)+rand(n,n)*i \end{array}$

When n = 100, the numerical comparison results are given in Fig. 4. As it is seen, the SGI-FOM(10) and SGI-GMRES(10) algorithms require fewer iteration steps than the the extended CGNE and GI methods for achieving the prescribed stopping criterion.



Figure 4: comparison results of Example 4.4

5. CONCLUSION

In this paper, we propose the SGI-FOM(m) method and SGI-GMRES(m) method to solve the Stein-like matrix equation. From numerical examples we conclude that the SGI-FOM(m) method and SGI-GMRES(m) method may be a useful tool for solving the Stein-like matrix equation. In the future, we will study the possibility of combining the augment technology with SGI-FOM(m) method and SGI-GMRES(m) method.

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